# PASSWORD CRACKING BASED ON LEARNED PATTERNS FROM DISCLOSED PASSWORDS 

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#### Abstract

Password-based authentication systems are still the most commonly used mechanism for protecting sensitive information despite being vulnerable to dictionary based attacks. To guard against such attacks, many organizations enforce complicated password-creation rules and require that passwords include numeric and special characters. This study demonstrates that as long as passwords are not difficult to remember, they remain vulnerable to "smart dictionary" attacks. In this study, a password analysis platform is developed to formally analyze commonly used passwords and identify frequently used password patterns and their associated probabilities. Based upon these patterns, we establish a model consisting of a Training set, a Dictionary set and a Testing set (TDT model) to generate probabilistic passwords sorted in decreasing order. The model can be used to dramatically reduce the size of the password space to be searched. Simulation results show that the number of passwords cracked using the TDT model is 1.43 and 2.5 times higher compared with the John-the-Ripper attack and Brute-force attack, respectively. We also design a hybrid password cracking system combining different attacks to verify the effectiveness of the proposed method. After applying the TDT model, the number of passwords cracked increased by up to 273\%.


Keywords: Password cracking, Dictionary attack, Brute-force attack, TDT model

1. Introduction. Identity authentication or access control through passwords is still a widely used method of ensuring system security, despite the increased use of alternative techniques such as graphical passwords [1], smart card, and biometrics. However, password based security is vulnerable to the dictionary attack [2]. In order to force users to create strong passwords, system administration policies often enforce several complex rules intended to force users into creating strong passwords [3,4]. Under such rules, users may be required to use numeric or special characters, have to enter passwords of a minimum length, and avoid words found in a dictionary. Users often struggle to create passwords that meet these requirements. However, it also becomes difficult to crack such passwords using a dictionary attack.

For example, consider the password 'computer123!'. If password rules were not enforced, a user may simply have selected 'computer' as their password. However, in order to satisfy the requirements of the system administration policy, they may have to append numeric characters ' 123 ' and a special character '!'. In this way, security is enhanced because the difficulty of cracking the password increases from $26^{8}$ to $68^{12}$.

In addition to enforcing high password complexity, many password-based cryptosystems execute multiple iterations in their core encryption algorithm to defend against the Bruteforce attack. In other words, they execute the core encryption algorithm many times to increase password complexity and cracking time. Take the WORD software as an example, where an attacker can crack more than 30,000 passwords per second for the 2003 edition using a dual quad-core 2.33 GHz Intel computer. However, it is only able to crack around 400 passwords per second for the 2007 edition. Therefore, taking WORD 2007 as an example, for a lowercase string of length 8 , it may take 16.5 years $\left(26^{8} / 400 * 86400 * 365\right)$ to crack using the Brute-force attack.

These examples suggest that due to rapidly changing user habits for password selection and enhancements in the encryption algorithm, the dictionary attack and brute-force attack face an enormous challenge. This paper, looking from the attacker's standpoint, aims to overcome this problem and design a good method for generating highly efficient passwords for cracking. The proposed method makes the dictionary attack stronger by overcoming its shortcomings and also overcomes the limitations of the brute-force attack.

In this paper, we analyze user habits in creating passwords, and design an effective password attack method. First, we design a password analysis platform using Flex [5,6] with Cygwin [7]. It uses regular expressions to categorize and analyze disclosed passwords to understand how users create passwords. By analyzing the disclosed passwords, we are able to determine how users incorporate the required character types in their passwords and summarize some frequently used password patterns with their associated probabilities. We then create a model comprising a Training set, a Dictionary set and a Testing set, called the TDT model. This model can generate passwords and sort them in decreasing order of probability. Finally, we devise a hybrid password cracking system, which is a three stage sequence comprising a dictionary attack, the TDT model attack, and a brute force attack. We then simulate and apply this system to cracking UNIX access passwords.

This paper is organized as follows. Section 2 describes related research on password attacks. Section 3 proposes pattern analysis from disclosed passwords. Section 4 analyses the strategy for generating passwords. Section 5 presents the simulation results. Section 6 demonstrates the effectiveness and comparisons with other related researches. Section 7 contains conclusion and future work.
2. Related Work. Techniques for cracking or acquiring passwords include social engineering, phishing and shoulder surfing. However, most current identity authentication attacks (focusing on passwords) are still based on the dictionary attack or on brute-force methods. Other attacks such as time-memory tradeoff [8,9] are gradually gaining importance as computing power and storage space increases. However, this attack requires substantial computational time and effort to build Rainbow tables [10] and thus has not been widely deployed. In addition, an effective defense has been mounted against Rainbow tables by exploiting the fact that some hash functions are combined with a random salt (for example salted SHA-1), which causes the same passwords to produce different hash values.

Actually, many companies such as Elcomsoft, Passware, and Wwwhack have developed password recovery software for documents files that use character strings for data encryption, such as Word, Excel, PDF, RAR, and ZIP encryption files. However, most of these password recovery packages are still based on the dictionary attack or brute-force attack.

The effectiveness of dictionary attack and brute-force attack has been challenged by the increasing awareness and emphasis on secure password policies. In addition, many password-based cryptosystems have incorporated more complicated protection mechanisms into existing encryption algorithms to defend against a brute-force attack, increasing the time required for every password guess in a brute-force attack. The MD5 crypt is one such example [11]. The MD5 hash can be brute-forced in a reasonable amount of time. However, applying the MD5 hash a thousand times increases the times required to crack the hash by the same factor. To overcome this protection method, password crackers have to design a good guessing mechanism that can guess effectively and obtain results in a limited amount of time.

Castelluccia [12] and Narayanan [13] proposed a password-cracking technique based on a Markov model, in which password guesses are based on the contextual frequency of characters. Weir et al. [14,15] presented a novel password-cracking technique that used the text structure from training data while applying mangling rules to the text itself. The authors found their technique to be more effective than the John the Ripper tool [16]. In a separate study, Zhang et al. [17] found Weir's algorithm to be the most effective among the techniques they used. In 2011, Kelley et al. [18] applied Weir's algorithm and a variation of the Markov model to generate simulated passwords and evaluate their strength.

However, Weir's algorithm inherently depends on common dictionaries. Weir et al. used three different dictionaries to generate passwords. Their method will not be effective if users have used words that are not in these dictionaries as passwords. For example, if the password is '!qazxsw@', based on the meaningless sequence of characters 'qazxsw', then Weir's algorithm will not be able to crack the password. Although 'qazxsw' is not a meaningful word, it is a keyboard pattern. This paper applies a variation of Weir's algorithm to generate more passwords more efficiently. We analyzed the patterns found in disclosed passwords, and developed the TDT model to generate passwords. These generated passwords were then used to crack UNIX access passwords.
3. Pattern Analysis from Disclosed Passwords. User's password usually comes from printable characters. These characters are used in forming the passwords. We define the elements and password patterns obtained from the disclosed passwords.
3.1. Basic elements and patterns of disclosed passwords. The basic elements of user's disclosed passwords are the printable characters. These characters available on a keyboard can be categorized into four different types, i.e., numeric $(N)$, lowercase $(L)$, uppercase $(U)$ and other $(O)$. Table 1 shows all the 94 printable characters and their type.

Table 1. Basic elements and types of users' passwords

| Type | Character Number | Basic Elements |
| :---: | :---: | :---: |
| Numeric $(N)$ | 10 | 0123456789 |
| Lowercase $(L)$ | 26 | abcdefghijklmnopqrstuvwsyz |
| Uppercase $(U)$ | 26 | ABCDEFGHIJKLMNOPQRSTUVWXYZ |
| Other $(O)$ | 32 | $\sim{ }^{\prime}$ @ $\# \$ \wedge^{\wedge} \&^{*}()_{-}-+=[] \backslash\{ \} \mid ; ;^{\prime}, . /<>?$ |

In order to describe and analyze the disclosed passwords comprising of $N, L, U$ and $O$ strings, a terminology is defined as follows.

Definition 3.1. ( $N^{+}, N_{n}$ ) A Number string (denoted $N^{+}$) is a continuous sequence of characters of $N$ type. A Number string of length $n$ (denoted $N_{n}$ ) is a continuous sequence of characters of $N$ type of length $n$.

Definition 3.2. (Element) An element of type $N_{n}$ is an instance of Number string $N_{n}$ found in a disclosed password.

For example, from an instance of disclosed passwords like '12348888', we obtain an element '12348888' of type $N_{8}$, which also belong to $N^{+}$. By Definition 3.1, $N^{+}$represents either one of $N_{1}, N_{2}, N_{3}, \ldots, N_{i}$ and denotes a super class among $N_{1}, N_{2}, N_{3}, \ldots, N_{i}$.

Lemma 3.1. The same definition also applies to Lowercase string $\left(L^{+}, L_{n}\right)$, Uppercase string $\left(U^{+}, U_{n}\right)$ and other string $\left(O^{+}, O_{n}\right)$. An Alpha string $\left(A^{+}, A_{n}\right)$ is a continuous sequence of characters of $L$ or $U$ types.

Users create their passwords as patterns based on a combination among the strings of character types, defined as follows:

Definition 3.3. (Password pattern and pattern class) The password pattern of a password is the combination of strings of character type ( $N, L, U$ and $O$ ) in a password. It is represented in terms of the combinations strings of $N_{n}, L_{n}, U_{n}$ and $O_{n}$, where the subscript $n$ represents the length of the corresponding type. The pattern class represents a set of password patterns and is a combination of strings of $N^{+}, L^{+}, U^{+}$and $O^{+}$.

For example, the password pattern of the password 'Password10!!' is ' $U_{1} L_{7} N_{2} O_{2}$ ' since the password consists of four elements, i.e., 'P', 'assword', '10', and '!!'. These elements belong to patterns $U_{1}, L_{7}, N_{2}, O_{2}$, respectively. Furthermore, the password pattern ' $U_{1} L_{7} N_{2} O_{2}$ ' is an instance of pattern class ' $U^{+} L^{+} N^{+} O^{+1}$.

Some examples are presented in Table 2. We can see that most passwords are based on Alpha strings; they consist of Lowercase or Uppercase strings. Number strings or Other strings of variable length are inserted before or appended after the Alpha strings. Some Alpha strings are dictionary words, such as 'password', 'computer', 'tiger', and 'count'. Other passwords are based on the keyboard layout, e.g., 'qwerty', 'qazwsx'. In Section 3.2, we develop a password analysis platform to analyze users' disclosed passwords in detail.

Table 2. Some instances from users' disclosed passwords

| Instance | Pattern Class | Password Pattern | Elements |
| :---: | :---: | :---: | :---: |
| password1 | $L^{+} N^{+}$ | $L_{8} N_{1}$ | password, 1 |
| computer88 | $L^{+} N^{+}$ | $L_{8} N_{2}$ | computer, 88 |
| 13tmein | $N^{+} L^{+}$ | $N_{2} L_{5}$ | 13, tmein |
| jester6\# | $L^{+} N^{+} O^{+}$ | $L_{6} N_{1} O_{1}$ | jester, $6, \#$ |
| Count!!!! | $U^{+} L^{+} O^{+}$ | $U_{1} L_{4} O_{4}$ | C, ount, !!!! |
| qwerty2009 | $L^{+} N^{+}$ | $L_{6} N_{4}$ | qwerty, 2009 |
| AAaa1211\&\& | $U^{+} L^{+} N^{+} O^{+}$ | $U_{2} L_{2} N_{4} O_{2}$ |  |
| @merican | $O^{+} L^{+}$ | $O_{1} L_{7}$ | @, merican |
| 95celica | $N^{+} L^{+}$ | $N_{2} L_{6}$ | 95, celica |
| 8qazwsx9 | $N^{+} L^{+} N^{+}$ | $N_{1} L_{6} N_{1}$ | 8, qazwsx, 9 |

3.2. Statistics of password patterns. In October 2008, a phishing attack was launched to obtain the email addresses and passwords of Rockyou users [19]. The phishing attack received much publicity, and the $14,344,389$ passwords obtained were made available for download. However, the downloaded password corpus did not contain any information about associated Rockyou usernames. In the paper, we use these passwords as disclosed real password data for further analysis.

To analyze the password corpus more systematically, we developed a password analysis platform using Flex [5,6] with Cygwin [7]. The Cygwin package provides a complete UNIX environment from within a Windows PC. After installing Cygwin, Flex will be able to complete all flex assignments on a PC. Flex is normally used to partition a stream of characters into tokens. It takes a specification that associates regular expressions with actions as input. By using this password analysis platform with regular expressions, we analyze the passwords from Rockyou as described next.

## Pattern Class, Password Pattern, and Element

According to our statistics, there are 141,278 pattern classes in Rockyou. However, 160 of these pattern classes are composed of continuous sequences of one to four character types, which cover $97.51 \%$ of these passwords in RockYou. For example, the password class of 'abc123' is ' $L^{+} N^{+1}$, consisting of two character types. The reason is that overly complex passwords are difficult for users to remember. Therefore, we focus on calculating possible combinations of continuous sequences of one to four character types.

With our password analysis platform, we represent these 160 combinations by regular expressions, and use Flex to obtain statistics. Table 3 lists the top fifteen pattern classes in Rock You disclosed passwords.

Table 3. Top fifteen pattern classes from Rockyou

| Pattern Class | Occurrences | Percentage of Total |
| :---: | :---: | :---: |
| $L^{+} N^{+}$ | $4,720,184$ | $32.91 \%$ |
| $L^{+}$ | $3,726,129$ | $25.98 \%$ |
| $N^{+}$ | $2,346,744$ | $16.36 \%$ |
| $N^{+} L^{+}$ | 499,167 | $3.48 \%$ |
| $L^{+} N^{+} L^{+}$ | 388,157 | $2.71 \%$ |
| $U^{+} N^{+}$ | 325,941 | $2.27 \%$ |
| $U^{+} L^{+} N^{+}$ | 236,331 | $1.65 \%$ |
| $U^{+}$ | 229,875 | $1.6 \%$ |
| $L^{+} O^{+} L^{+}$ | 172,279 | $1.2 \%$ |
| $L^{+} O^{+} N^{+}$ | 144,129 | $1.0 \%$ |
| $L^{+} N^{+} L^{+} N^{+}$ | 123,344 | $0.86 \%$ |
| $L^{+} O^{+}$ | 122,560 | $0.85 \%$ |
| $N^{+} L^{+} N^{+}$ | 118,877 | $0.83 \%$ |
| $U^{+} L^{+}$ | 98,515 | $0.69 \%$ |
| $L^{+} N^{+} O^{+}$ | 65,164 | $0.45 \%$ |

From Table 3, the top three pattern classes account for $75.25 \%$ of all passwords. This indicates that most users use a combination of Lowercase and Number strings, or just Lowercase strings or Number strings as their passwords. The top ranked ' $L^{+} N^{+\prime}$ pattern class comprises $32.91 \%$ of the total. The statistical information regarding password class gives us a more detailed look into how users incorporate the four character types into passwords. For the ${ }^{\prime} L^{+} N^{+1}$ pattern class, we further calculate the password patterns according to various lengths of the Lowercase and Number strings. The result is shown
in Table 4. The most popular password pattern is ' $L_{6} N_{2}$ ', consisting of Lowercase string of six characters followed by Number string of two characters. The probability assigned to each password pattern is

$$
\operatorname{Probability}(x)=\frac{x_{\mathrm{num}}}{K_{\mathrm{tot}}}
$$

where $x_{\text {num }}$ is the number of occurrences of that particular password pattern $x$, and $K_{\text {tot }}$ is the total number of ' $L^{+} N^{+\prime}$ pattern class. For instance, the probability of password pattern ' $L_{6} N_{2}$ ' is $420318 / 4720184=8.91 \%$. The same statistical information is recorded for each of the pattern class.

After obtaining the statistics of password patterns for each of the pattern class, we can sort the password patterns according to their probabilities. Next, we parse the passwords according to the previously defined four strings ( $N, L, U, O$ ) and then count their probabilities respectively.

Table 4. Top ten password patterns in ${ }^{\prime} L^{+} N^{+\prime}$ pattern class from Rockyou

| Password Patterns | Occurrences | Percentage of ' $L^{+} N^{+1}$ |
| :---: | :---: | :---: |
| $L_{6} N_{2}$ | 420,318 | $8.91 \%$ |
| $L_{5} N_{2}$ | 292,306 | $6.19 \%$ |
| $L_{7} N_{2}$ | 273,624 | $5.80 \%$ |
| $L_{4} N_{4}$ | 235,360 | $4.99 \%$ |
| $L_{4} N_{2}$ | 215,074 | $4.56 \%$ |
| $L_{8} N_{2}$ | 213,109 | $4.51 \%$ |
| $L_{6} N_{1}$ | 193,097 | $4.10 \%$ |
| $L_{7} N_{1}$ | 189,847 | $4.02 \%$ |
| $L_{5} N_{4}$ | 173,559 | $3.68 \%$ |
| $L_{6} N_{4}$ | 160,592 | $3.40 \%$ |

## Number strings in Passwords

The selection of Number strings in passwords often represents something of significance, such as birthday, telephone number or special date. The patterns of Number strings are been classified dependent upon their lengths in passwords. Table 5 describes the statistics of length between one and ten in Number strings. The pattern of length two is the most frequent as its occurrence is $24.38 \%$. $65.35 \%$ of these patterns have lengths between one and four for Number strings. Each pattern of Number strings is analyzed to determine which elements are most frequently seen in passwords. Table 6 shows the top ten elements of length two of pattern ' $N_{2}$ '. The most frequent element is ' 12 ', with an occurrence of 4.81\%. Similar statistics are recorded for each of pattern in Number strings.

## Other strings in Passwords

Other strings are also important in password generation. The same process that is detailed above for Number strings is repeated for Other strings. The patterns of length between one and ten in Other strings occur a total of 301,096 times. Table 7 displays the frequencies of each pattern of length between one and ten for Other Strings. It is clear that pattern of length one occurs more frequently than any others do. Table 8 describes the most frequent element of length one occurring in pattern ' $O_{1}$ '. The same statistics are collected for the other patterns of Other strings and recorded.

## Alpha Strings in Passwords

Alpha strings are the main components in creating a password, which are comprised of Lowercase strings and Uppercase strings. In order to enhance the strength of passwords, usually the passwords formed by familiar Alpha strings are combined with Number strings

Table 5. Patterns of length between one and ten in Number strings

| Patterns of Number Strings | Occurrences | Percentage of Total |
| :---: | :---: | :---: |
| $N_{2}$ | $2,131,337$ | $24.38 \%$ |
| $N_{4}$ | $1,330,610$ | $15.22 \%$ |
| $N_{1}$ | $1,321,553$ | $15.12 \%$ |
| $N_{3}$ | 928,760 | $10.63 \%$ |
| $N_{6}$ | 776,436 | $8.88 \%$ |
| $N_{7}$ | 608,993 | $6.97 \%$ |
| $N_{8}$ | 523,130 | $5.98 \%$ |
| $N_{10}$ | 503,930 | $5.77 \%$ |
| $N_{9}$ | 343,396 | $3.93 \%$ |
| $N_{5}$ | 272,088 | $3.11 \%$ |

Table 6. Elements of length two of pattern ' $N N_{2}$ ' in Number strings

| Elements of Length Two | Occurrences | Percentage of ' $N_{2}{ }^{\prime}$ |
| :---: | :---: | :---: |
| 12 | 102,590 | $4.81 \%$ |
| 13 | 76,775 | $3.60 \%$ |
| 11 | 65,201 | $3.06 \%$ |
| 22 | 58,058 | $2.72 \%$ |
| 23 | 57,825 | $2.71 \%$ |
| 07 | 55,693 | $2.61 \%$ |
| 21 | 55,192 | $2.59 \%$ |
| 01 | 53,608 | $2.52 \%$ |
| 10 | 53,510 | $2.51 \%$ |
| 14 | 53,300 | $2.50 \%$ |

Table 7. Patterns of length between one and ten in Other strings

| Patterns of Other Strings | Occurrences | Percentage of Total |
| :---: | :---: | :---: |
| $O_{1}$ | 238,652 | $79.26 \%$ |
| $O_{2}$ | 38,780 | $12.88 \%$ |
| $O_{3}$ | 16,184 | $5.37 \%$ |
| $O_{4}$ | 3,544 | $1.18 \%$ |
| $O_{6}$ | 1,260 | $0.42 \%$ |
| $O_{5}$ | 1,177 | $0.38 \%$ |
| $O_{7}$ | 623 | $0.21 \%$ |
| $O_{8}$ | 452 | $0.15 \%$ |
| $O_{9}$ | 226 | $0.08 \%$ |
| $O_{10}$ | 198 | $0.07 \%$ |

and Other strings. For example, password 'computer123', consists of the Alpha string 'computer' and Number string '123'. In this paper, we take Alpha strings as the basis and associate them with Number strings and Other strings to generate new passwords. Therefore, we parse out the Alpha strings from Rockyou and use the set as the Dictionary set, which is illustrated in Section 4. The total number of unique Alpha strings from Rockyou is 5,521,967.

According to our observation for these parsed Alpha stings, many users still select familiar dictionary words. In order to count the number of dictionary words in Alpha

Table 8. Elements of length one in pattern ' $O_{1}$ '

| Elements of Length One | Occurrences | Percentage of ' $O_{1}{ }^{\prime}$ |
| :---: | :---: | :---: |
| $!$ | 64,602 | $27.07 \%$ |
| $\cdot$ | 42,446 | $17.79 \%$ |
|  | 33,950 | $14.23 \%$ |
| $\$$ | 13,995 | $5.86 \%$ |
| - | 10,941 | $4.58 \%$ |
| - | 8,393 | $3.52 \%$ |
| $?$ | 6,922 | $2.90 \%$ |
| $($ | 6,447 | $2.70 \%$ |
| ) | 6,085 | $2.55 \%$ |
|  | 4,818 | $2.02 \%$ |

Table 9. Statistics of dictionary words in alpha strings in Rockyou

| Alpha strings | Numbers of Occurrences | Percentage of Total |
| :---: | :---: | :---: |
| Dictionary words | $1,254,021$ | $8.74 \%$ |

strings, we collected 869,228 dictionary words from Dic-0294 [20] for reference. To enhance the efficiency of comparison, the MD5 hashing algorithm is used to process the large number of Alpha strings. Tests show this provides a ten-time speedup. The statistics are given in Table 9. It proved that many users still select weak dictionary words to form their passwords.

## Keyboard Strings in Passwords

In this paper, we also analyze the password patterns generated by the keyboard positions. The Keyboard strings refer to passwords formed by relative keyboard positions and associated finger movements. Its examples include 'keyboard neighbor' and 'keyboard parallel $n^{\prime}$, which are parallel patterns of length $n$. In this paper, such passwords are referred to as Keyboard strings, and represented by the symbol $K$. For example, the Keyboard string '1qa2ws' is formed by a keyboard parallel 3 pattern on the keyboard. Using Flex, we overcome the limitation on recursive expressions, and use regular expressions to analyze keyboard neighbor and parallel relationships. The statistics are given in Table 10. Notice that although the percentage of keyboard strings is not high, with increasing awareness and emphasis on password security, more and more users choose this structure for their passwords.

Table 10. Statistics of keyboard strings in Rockyou

| Keyboard String | Example | Numbers of Occurrences | Percentage of Total |
| :---: | :---: | :---: | :---: |
| Keyboard Neighbor | asdfvcxz | 20,739 | $0.15 \%$ |
| Keyboard Parallel 3 | 1qa2ws | 1898 | $0.013 \%$ |
| Keyboard Parallel 4 | 0okm9ijn | 780 | $0.0054 \%$ |
| Keyboard Parallel 5 | 12345qwert | 231 | $0.0016 \%$ |

Using our password analysis platform, we analyzed the passwords from Rockyou and found 176 common and popular password patterns through Flex. These patterns cover $96.3 \%$ of the passwords in Rock You corpus. With this analysis platform, the patterns of users' disclosed passwords can be analyzed quickly. Moreover, this platform is also easy to maintain and extend as password patterns can be added and removed. For instance, if in the future we find more than four character types change used in passwords, e.g.,
'!!!Password88\#\#', containing ' $O_{3} U_{1} L_{7} N_{2} O_{2}$ ', i.e., a continuous sequence of five character types, we just need to add the regular expression that describes the new pattern to Flex before performing analysis. In Section 4, we use these statistics of the password patterns as a basis to generate efficient passwords for reducing the number of guesses and increasing the hit rate.
4. Strategy for Generating Passwords. From statistics of disclosed passwords for RockYou described in Section 3, we summarize some common patterns for selecting passwords.

1) Many passwords are composed of Lowercase and Number strings, i.e., ' $L^{+} N^{+}$', the Lowercase strings are likely appended by Number strings.
2) Number or Other strings are inserted before or appended after Alpha strings, e.g., 'jessica!!' belongs to ' $L_{8} O_{2}$ ' pattern, the Alpha string 'jessica' is appended with the Other string '!!'.
3) Alpha strings include Keyboard strings, e.g., 'qwerty123' belongs to ' $K_{6} N_{3}$ ' pattern, containing the Keyboard string 'qwerty'.
4) Alpha strings include dictionary words, e.g., 'football1234' belongs to ' $L_{8} N_{4}$ ' pattern, containing the dictionary word 'football'.
Instead of trying to create rules that mimic these common password patterns, we assign probabilities to these patterns, and then use those probabilities directly to generate very fine-grained rules, and thus generate passwords for cracking.
4.1. TDT model. Based on the analysis of the disclosed passwords in Section 3, we create the TDT model consisting of the Training set, Dictionary set and Testing set, to generate the passwords. Three sets in the model are defined as follows.

Definition 4.1. (Training set, Dictionary set and Testing set) The Training set specifies users' disclosed passwords. The Dictionary set specifies dictionary words, Alpha strings or Keyboard strings. The Testing set specifies the cracking targets, which are plaintext passwords. If a password is used in the Training set, we ensure that it is not included in the Testing set.

The TDT model uses the Training set to create password rules, and use these rules along with the Dictionary set to generate passwords. Then it uses the generated passwords to crack passwords in the Testing set. Note that a password is considered 'cracked' if our method generates a guess that matches a password in the Testing set. The framework of the TDT model is shown in Figure 1 and is described below.

1) Disclosed passwords collection: In our method, user disclosed passwords are used as the Training set, collected from various sources (e.g., specific websites) or from broken passwords of target. For this study, we simply use the Rockyou password corpus as the Training set.
2) Establishing UDP rules: When sufficient passwords have been collected, the Training set is used to create user disclosed passwords rules (shortened as UDP rules).
3) Passwords generation: The passwords could be generated using the UDP-rules along with the Dictionary set. The Dictionary set consists of the commonly used dictionary words or Alpha and Keyboard strings parsed from the Training set.
4) Password cracking and feedback: The generated passwords are used to crack passwords from the Testing set. Once passwords have been cracked, they can be fed back to the Training set. When the Training set is updated, we repeat steps $2-4$, to update UDP rules and generate new passwords. Through this repeated learning and training


Figure 1. The framework of TDT model
process, we can optimize the accuracy of the UDP rules and enhance the guessing ability of the model.
4.2. Establishing UDP rules. Each password in the Training set is parsed. The parsed strings in a password that have Number strings and Other strings are called 'leaf nodes' and those having Lowercase, Uppercase, and Keyboard strings are called 'inner nodes'. Leaf nodes are categorized based on character types and their length. Their probabilities are calculated based on the number of times they occur in the Training set. The probability of each leaf node is calculated as follows.

## $\frac{\text { number of occurrences of specific leaf node of length } k}{n_{k}}$

where $n_{k}$ is the total number of strings of length $k$. For example, if the Number strings of length 3 are ' 123 ' and ' 456 ', and they occur in the Training set 100 and 60 times respectively, then the probability for string ' 123 ' is $100 / 160=0.625$ and that for string ' 456 ' is $60 / 160=0.375$. Inner nodes are used as dictionary words and added to the Dictionary set. At this point, dictionary words in the Dictionary set do not have probabilities associated with them, so we set the probability for inner nodes is 1.0 .

Moreover, we use the symbols $N_{n}, L_{n}, U_{n}, O_{n}, K_{n}$ defined in Section 3 to represent each password pattern, where $N$ stands for Number, $L$ for Lowercase, $U$ for Uppercase, $O$ for Other, and $K$ for Keyboard, and the subscript $n$ represents the length of the corresponding symbols. For example, the password pattern of 'abc123' is ' $L_{3} N_{3}$ '. Based on this rule, passwords with the same patterns will be merged, and the probability of each pattern can be calculated. Thus, we can calculate the probability of each password pattern and find which among them are used more frequently.

We describe this process through a simple example. Assume that there are eight passwords in the Training set as shown in Figure 2. For leaf nodes, we categorize them by length, and then calculate their probabilities. Among strings of length $=1$, the Other string '@' appears twice, and '!' appears once. Therefore, the probability of '@' is $2 / 3=0.67$, while that of '!' is $1 / 3=0.33$. Using the same method, we can find the probabilities of Number strings of length three and four, as shown in Figure 2(a). Inner nodes are used as dictionary words, as explained in Section 4.3. In addition, after converting the eight passwords into their patterns, we merge these patterns and obtain five distinct patterns, as shown in Figure 2(b). For example, the three passwords 'abc123',
'cat789' and 'man666' all have the ' $L_{3} N_{3}$ ' pattern, so its probability is $3 / 8=0.375$. Note that the higher the probability of a pattern, the more frequently it is used by users.
4.3. Generating passwords. Using the UDP rules created in Section 4.2, we generate passwords with the help of a Dictionary set. The Dictionary set mainly consists of commonly used dictionary words, or the inner nodes parsed from the Training set. By constructing the Dictionary set, we can gain a larger coverage of the target set of passwords. The core idea of this paper is to generate highly efficient passwords by using the Dictionary set as the basis and combining it with rules generated from the Training set. An example Dictionary set is shown in Figure 3. In addition to the eight inner nodes parsed from the Training set in Section 4.2, we added the three words 'princess', 'nicole' and 'you' to the Dictionary set.
(a)


Figure 2. Examples of generated UDP rules using a training set

| Training Set |  |
| :--- | :--- |
| abc123 <br> password456 <br> !qwerty@ <br> cat789 <br> man666 <br> computer456 <br> football1970 <br> Hero@ | Parsing |
| Dictionary Set <br> words | abc <br> password <br> qwerty <br> cat <br> man <br> computer <br> football <br> Hero |

Figure 3. The dictionary set

Next, we use the Dictionary set and the UDP rules to generate passwords and calculate their probabilities. The formal mathematical definition is as follows.

Theorem 4.1. Assume that $P\left(\beta_{x}\right)$ be the probability of the pattern of a password $x$, which $\beta_{x}=t_{1} t_{2} t_{3} \ldots t_{n}$ is the pattern of a password $x$, consisting of $t_{1}, t_{2}, t_{3}, \ldots, t_{n}$ inner or leaf nodes. Let $v_{\ell_{i}}\left(t_{i}\right)$ be the probability of the $i$-th inner or leaf nodes of length $\ell$ in the Training set. Therefore, the probability of generated password $x$ is given by

$$
P(x)=P\left(\beta_{x}\right) \cdot \prod_{i=1}^{n} v_{\ell_{i}}\left(t_{i}\right), t_{i} \in\{\text { inner nodes, leaf nodes }\}
$$

We use the generated password 'abc456' as an example. Since the string 'abc' in the Dictionary set matches ' $L_{3}$ ' in the password pattern ' $L_{3} N_{3}$ ', it must be combined with Number strings in the form of ' $N_{3}$ ' to generate a password. Therefore, the probability of generated password 'abc456' is

$$
\begin{aligned}
P(a b c 456) & =P\left(\beta_{a b c 456}=t_{1} t_{2}=L_{3} N_{3}\right) \cdot v_{\ell_{1}}\left(L_{3}\right) \cdot v_{\ell_{2}}\left(N_{3}\right) \\
& =0.375 \cdot 1.0 \cdot 0.4 \\
& =0.15
\end{aligned}
$$

For the same method, we calculate the probability of each password based on the respective probabilities of Number strings '123' '789' and '666', as shown in Table 11.

Table 11. Probability of generated passwords using dictionary set and UDP rules

| Dictionary set | Pattern | Number String | Probability of <br> Generated Passwords | Generated <br> Passwords |
| :---: | :---: | :---: | :---: | :---: |
|  | Probability | Probability | $0.375 * 0.4=0.15$ | abc456 |
| abc | $L_{3} N_{3}$ | 456 |  |  |
|  | 0.375 | 0.4 | $0.375 * 0.2=0.075$ | abc123 |
| abc | $L_{3} N_{3}$ | 123 |  |  |
|  | 0.375 | 0.2 | $0.375 * 0.2=0.075$ | abc789 |
| abc | $L_{3} N_{3}$ | 789 |  |  |
|  | 0.375 | abc666 |  |  |
| abc | $L_{3} N_{3}$ | 666 |  |  |
|  | 0.375 | 0.2 |  |  |

Since the probability of each generated password can be calculated, we sort the probabilities of generated passwords in decreasing order and apply a threshold $\theta$. Therefore, the generated passwords from the TDT model are

$$
D_{v, \theta}=\left\{X: P(\beta) \cdot \prod_{i=1}^{n} v_{l_{i}}\left(t_{i}\right) \geq \theta, t_{i} \in\{\text { inner nodes, leaf nodes }\}\right\}
$$

Table 12 shows the top twenty passwords along with their probabilities, as generated using the UDP rules in Section 4.2 along with our Dictionary set. The threshold $\theta$ is set as 0.075 . One of the advantages of our method is the ability to sort the generated passwords according to their probabilities. In this way, based on user requirements, we can set a probability threshold and filter out passwords having low probabilities, and thus reduce the guessing space. When a password is successfully cracked by our method, we can move the password back to the Training set, and rebuild the UDP rules. Through this repeat training mechanism, the generated passwords become better adapted to the targets, increasing the hit rate.

Table 12. Top 20 generated passwords using UDP rules and dictionary set

| Rank | Generated <br> passwords | Probability | Rank | Generated <br> passwords | Probability |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | abc456 | $0.375 * 0.4=0.15$ | 11 | password456 | $0.25 * 0.4=0.1$ |
| 2 | cat456 | $0.375 * 0.4=0.15$ | 12 | princess456 | $0.25 * 0.4=0.1$ |
| 3 | man456 | $0.375 * 0.4=0.15$ | 13 | Hero@ | $0.125 * 0.67=0.084$ |
| 4 | you456 | $0.375 * 0.4=0.15$ | 14 | abc666 | $0.375 * 0.2=0.075$ |
| 5 | computer1970 | $0.125 * 1.0=0.125$ | 15 | abc789 | $0.375 * 0.2=0.075$ |
| 6 | football1970 | $0.125 * 1.0=0.125$ | 16 | abc123 | $0.37 * 0.2=0.075$ |
| 7 | password1970 | $0.125 * 1.0=0.125$ | 17 | cat666 | $0.375 * 0.2=0.075$ |
| 8 | princess1970 | $0.125 * 1.0=0.125$ | 18 | cat789 | $0.375 * 0.2=0.075$ |
| 9 | computer456 | $0.25 * 0.4=0.1$ | 19 | cat123 | $0.375 * 0.2=0.075$ |
| 10 | football456 | $0.25 * 0.4=0.1$ | 20 | man666 | $0.375 * 0.2=0.075$ |

## 5. Experiments and Results.

5.1. Password generation and cracking. For a better computational complexity, we divided the $14,344,389$ leaked Rockyou passwords randomly into 28 subsets using the GNU shuf tool. Each such subset consisted of 500,000 passwords. For convenience, we selected the first and last subset as the Training set and Testing set respectively, called Rockyou1 and Rockyou28. The Dictionary set used is Dic-0294, which is a commonly used password cracking dictionary [20]. This dictionary contains 869,228 unique words. To evaluate performance, we selected Rockyou28, Myspace, and Phpbb plaintext passwords as the Testing set. The overall structure of TDT model is shown as Figure 4. The Myspace and Phpbb plaintext passwords are also created by attackers who had performed a phishing attack against users but failed to secure their collection server. Researchers proceeded to log into the malicious server and gathered the passwords from the hacker before the hacker could shut it down. The complete Myspace and Phpbb contained 36,764 and 184,380 plaintext passwords respectively. A summary is given in Table 13, and experiment results are presented in Figure 5.


Figure 4. Overall structure of TDT model
Figure 5 shows that our method is more efficient for password cracking. When the number of generated passwords is between 50 and 450 million, the hit rate for Myspace is $27-33 \%$, for Phpbb is $26-31 \%$, and for Rockyou 28 is $20-28 \%$. The results show that the hit rate for Rockyou28 is lower than that for Myspace and Phpbb. In other words, password strength in the Rockyou28 set is higher than that in Myspace and Phpbb. Password strength differs with users' cultural backgrounds, age, password creation requirements, importance of the password protected material, etc.

Table 13. The training, dictionary and testing sets

| Set | Name | Number |
| :---: | :---: | :---: |
| Training set | Rockyou1 | 500,000 |
| Dictionary set | Dic-0294 | 869,228 |
| Testing set | Rockyou28 | 500,000 |
|  | Myspace | 36,764 |
|  | Phpbb | 184,389 |



Figure 5. Percentage of passwords cracked

In Figure 5, we only use common used dictionary words (Dic-0294) as the basis for the Dictionary set. This procedure is suitable when we know nothing about the target passwords (i.e., the Testing set). However, we may have access to some passwords of specific targets from various channels. Having access to the partial list of target passwords, these passwords may be used to generate rules, and to parse the inner nodes from these passwords, combining them with the Dictionary set. In the above experiment, since Rockyou1 and Rockyou28 originate from the same source (Rockyou), we can assume they both have similar underlying password creation habits. Therefore, we consider Rockyou1 as the cracked passwords of Rockyou (i.e., the Training set), and consider Rockyou28 as the uncracked passwords of Rockyou (i.e., the Testing set). After parsing the inner nodes from Rockyou1 and combining them with the Dictionary set, the generated passwords should help increase the hit rate for Rockyou28.

To verify this prediction, we parsed out 284,903 inner nodes that included Lowercase, Uppercase, and Keyboard strings from Rockyou1, and call it Rockyou1_dic. These strings are added to Dic-0294 to form a new Dictionary set, which contained 1,154,131 dictionary words. The structure is shown in Figure 6 and the number of sets is listed in Table 14. We repeated the same experiments again and the experimental results are shown in Figure 7. When the number of generated passwords is between 50 and 450 million, the hit rate compared with Figure 5, increases from $27-33 \%$ to $30-36 \%$ for Myspace, from $26-31 \%$ to $29-34 \%$ for Phpbb, and from $20-28 \%$ to $32-41 \%$ for Rockyou28. It is apparent that the performance on Rockyou 28 shows the most improvement. This experiment shows that our UDP model can be customized for a specific target. When we add the inner nodes of Rockyou1 to the Dictionary set, the generated passwords have a higher chance of


Figure 6. Overall structure of TDT model
Table 14. The training, dictionary and testing sets

| Set | Name | Number |  |
| :---: | :---: | :---: | :---: |
| Training set | Rockyou1 | 500,000 |  |
| Dictionary set | Dic-0294 | 869,228 |  |
|  | Rockyou1_dic | 284,903 |  |
|  | $1,154,131$ |  |  |
| Testing set | Rockyou28 | 500,000 |  |
|  | Myspace | 36,764 |  |
|  | Phpbb | 184,389 |  |



Figure 7. Percentage of passwords cracked. The Dictionary set consists of Dic-0294 and Rockyou1_dic.
matching passwords in Rockyou28. This implies that if we know some password patterns of a specific target, the hit rate of password cracking can be improved.

To summarize the experimental results from Section 5.1, our method is suitable for two situations, described as follows:

- "Generalization" condition: where nothing is known about the target passwords. The Training set includes disclosed passwords (such as from Rockyou1) and the Dictionary set only includes commonly used dictionary words (e.g., Dic-0294). Figure 5 shows that using this information to crack passwords of Myspace, Phpbb and Rockyou 28 , a hit rate of $20 \%-33 \%$ can be achieved for 50 and 450 million generated passwords.
- "Customization" condition: where some information about the target passwords is known. The Training set includes the cracked target passwords (such as Rockyou1) and the Dictionary set includes the inner nodes parsed from these previously cracked passwords, as well as commonly used dictionary words (e.g., Dic-0294). Figure 7 shows that this method increases the hit rate of Rock28 from $20-28 \%$ to $32-41 \%$.
5.2. Experimental observations. Figure 8 shows a comparison of our method with the John-the-Ripper cracker tool [16] and Brute-force attack. The experiment used Rockyou28 as the Testing set. The result shows that the same number of generated passwords, our method cracked significantly more passwords than the John-the-Ripper and Brute-force attack. For example, for 300 million generated passwords, our method cracked $143 \%$ of the passwords cracked using John-the-Ripper, and $250 \%$ of those cracked using Brute-force attack.


Figure 8. Comparison with john-the-ripper and brute-force attack
Based on the above experimental results, the advantages of the proposed method are as follows:

1) The proposed method has self-training and learning capabilities. When more passwords of the target are known, more accurate passwords can be generated using the training and learning process.
2) Experiments showed that many users still use Alpha strings as the basis of their passwords, which are combined with variable-length Number or Other strings at the beginning or end of the Alpha strings.
3) In the past, people often used dictionary words for Alpha strings. However, Alpha strings among the disclosed passwords are increasingly keyboard words or other meaningless words. The proposed method is good at guessing such words as we add these words to the Dictionary set.
4) Based on the same number of generated passwords, our method can crack more passwords than the John-the-Ripper and Brute-force attack, as shown in Figure 8.
5) This method is designed for the English language at present. However, it can also be extended to other languages such as Chinese and Spanish.
6. Application. This section describes a hybrid password cracking system to crack access passwords collected from UNIX. Within this system, we introduce a new stage called TDT-model attack, which uses the passwords generated by TDT model.
6.1. Effectiveness. The design goal of the hybrid password cracking system is to crack encrypted target passwords by incrementally increasing the size of the password cracking space. Therefore, the system consists of three attack stages: the Dictionary attack followed by the TDT-model attack, and the Brute-force attack; this system is called the DTB password cracking system, as shown in Figure 9. The cracking strategy is to crack encrypted passwords using the Dictionary attack and the TDT-model attack first in finite time, and to then run the Brute-force attack when some encrypted passwords have not been cracked. The system can be used in DTB mode where all stages are enabled or in DB mode where only the TDT-model attack stage is disabled.

DTB Password Cracking System


Figure 9. The system architecture of DTB password cracking system

## UNIX Access Password Attack

We collected 382 encrypted access passwords for the UNIX system using the John-theRipper cracker tool [16]. These passwords are of 13-byte values generated by the DES cryptosystem. Then, we proceeded to crack the passwords in DTB and DB modes. The results showed that, in the same 5 -hour experiments, 226 passwords and 155 passwords are cracked in DTB mode and DB mode, respectively. After the Dictionary attack, in the same period of time, the number of passwords cracked with the TDT-model increases by 71 and by up to $273 \%[(75+37) / 41]$. The details are shown in Table 15, where the time taken to crack the passwords is given in brackets.

Table 15. Number of UNIX passwords cracked

| DTB Password <br> Cracking <br> System | DTB mode <br> Dictionary <br> attack | TDT-model <br> attack | Brute-force <br> attack | Dictionary <br> attack | Brute-force <br> attack |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | 114 <br> $(12 \mathrm{~min})$ | 75 <br> $(13 \mathrm{~min})$ | 37 <br> $(4 \mathrm{hr} 35 \mathrm{~min})$ | 114 <br> $(12 \mathrm{~min})$ | 41 <br> $(4 \mathrm{hr} 48 \mathrm{~min})$ |
|  | 226 |  |  | 155 |  |

6.2. Comparison. Among the dictionary-based password cracking methods, Weir et al. $[14,15]$ focused primarily on creating a probabilistic context-free grammar to generate word-mangling rules based on previously disclosed passwords. In 2011, Kelley et al. [18] proposed an efficient technique for evaluating password strength that is also based on the Weir's algorithm. However, Weir's algorithm will not be effective if users have not used passwords in their inherently common dictionaries. In the TDT model, we utilized Number, Lowercase, Uppercase, Other and Keyboard strings to generate rules. Moreover,
the TDT model may parse Lowercase, Uppercase, and Keyboard strings from the target's known passwords and combine these strings with the set of dictionary words, effectively generating "customized" passwords for the target. In addition, cracked passwords in the TDT model can be fed back into the Training set, and training rules can be reapplied. Table 16 showed a comparison of the methods of Weir, Kelley, and the TDT model.

Table 16. Comparisons of dictionary-based password cracking methods

| Method | Weir | Kelley | TDT model |
| :---: | :---: | :---: | :---: |
| Feature | Password structure <br> formally modeled using <br> a context-free grammar. | Measured password <br> strength by simulating <br> Weir's algorithms. | Developed the <br> probability-based password <br> cracking algorithm, a <br> superset of Weir's patterns; <br> verified in real cracking <br> system. |
| Execution | Offline. | Offline. | Offline/Online. |
| Flexibility | 'General' passwords <br> generated. | 'Specific' passwords <br> generated through different <br> composition policies. | 'Customized' passwords <br> can also be generated. |
| Limitation | Password generation <br> depend inherently on <br> common dictionaries. | Password strength depends <br> on the specific algorithm. | Password generation based <br> on known password <br> patterns. |

Figure 10 compares the cracking performance of the Weir method and our TDT model. The Myspace list containing 36,764 passwords used in the experiment was divided equally into a Training set and a Testing set. Dic-0294 was used as the Dictionary set. The performance of the TDT Model is clearly better than the Weir method.

Passwords generated using the TDT model are based on users' disclosed passwords and dictionary words, and can be used widely for cracking password-based cryptosystems. To resist attacks, many current password-based cryptosystems not only strengthen their cryptographic security, but also force users to create 'strong' passwords, which must at least be a combination of two character types and have a minimum length of eight characters. The strong passwords fall within the password space of the TDT model, and therefore, the cracking time is reduced.


Figure 10. Comparison between the Weir method and TDT model
7. Conclusion and Future Work. We developed a password analysis platform to formally analyze disclosed passwords, and establish the TDT model for generating passwords. Based on the probabilistic patterns of the TDT model, the passwords generated are sorted in decreasing order of probability; this is known as the TDT-model attack. We also design a hybrid password cracking system consisting of the Dictionary attack, TDT-model attack and Brute-force attack to verify the effectiveness of the TDT model. Experimental results show that this hybrid model is effective for cracking UNIX access passwords and there is a performance improvement of up to $273 \%$ from using the TDT-model attack. In future work, we plan to focus on letter replacement, such as replacing 'a' in a dictionary word with '@' to further strengthen passwords by providing more flexibility and generality. The effectiveness of this strategy with other languages may also be studied.

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